

# Hybrid Consensus: Efficient Consensus in the Permissionless Model

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IC3

# Blockchains

Satoshi Nakamoto, 2009

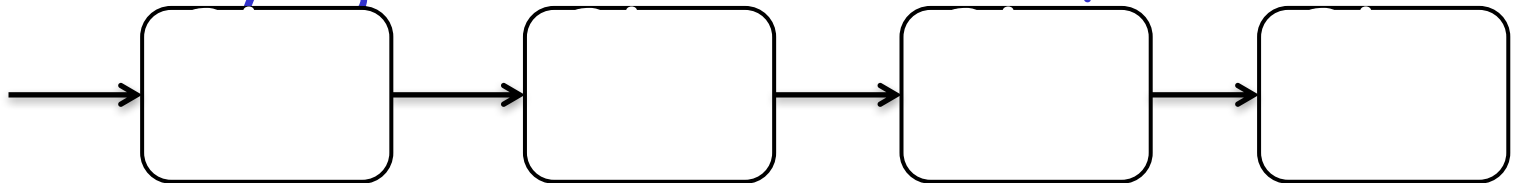
Public database

Shared & maintained between network participants



Blockchain agreement protocol

Periodically agree on a new block of data

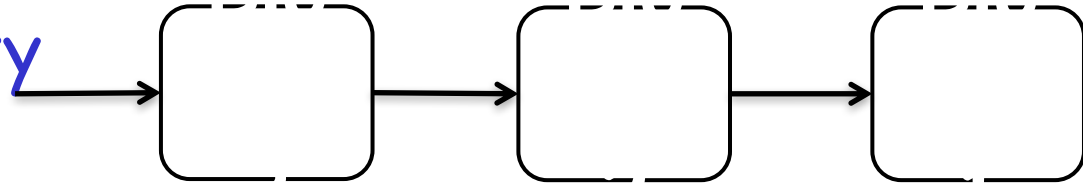


# The Blockchain Agreement Problem

## Set of $N$ computational nodes

## No inherent identity

## No PKI



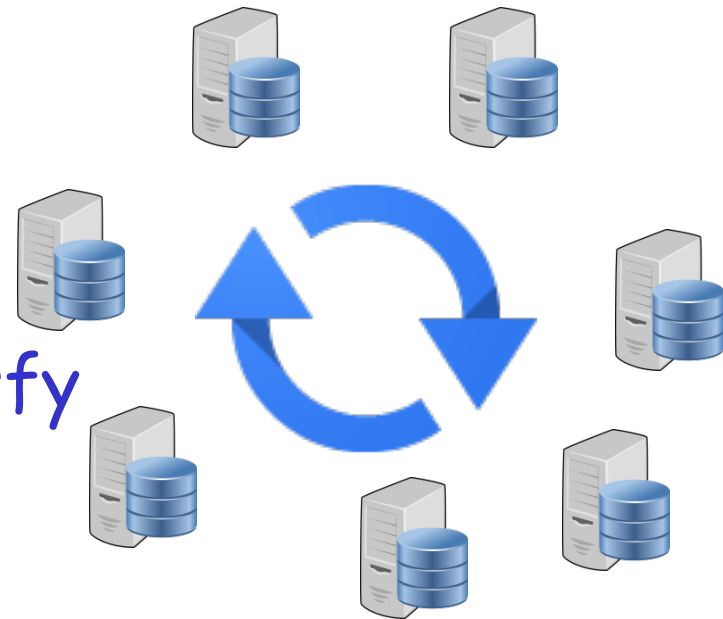
## Agree on a block

Includes a set of new transactions (or data)

## Two properties

# Agreement

Validity: all transactions satisfy some validity conditions



# Scalability Issue



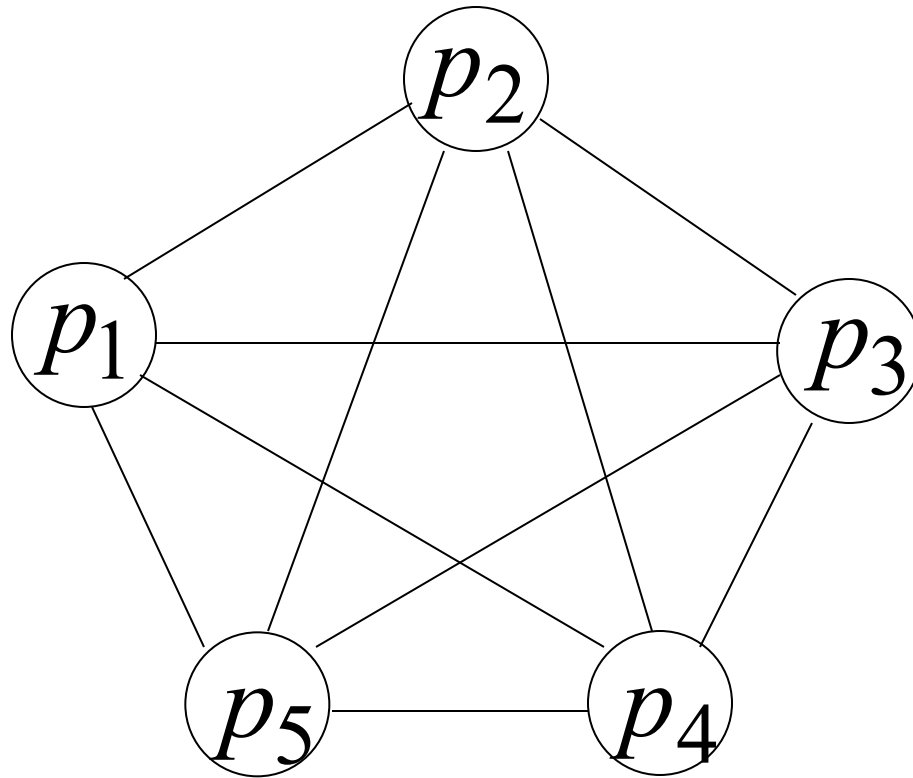
1 MB block per 10  
mins  
3-7 TXs per second

Support limited computations  
Several DoS attacks recently

Demand from Practice: 1,200 - 50,000 TXs/s



# Byzantine Agreement



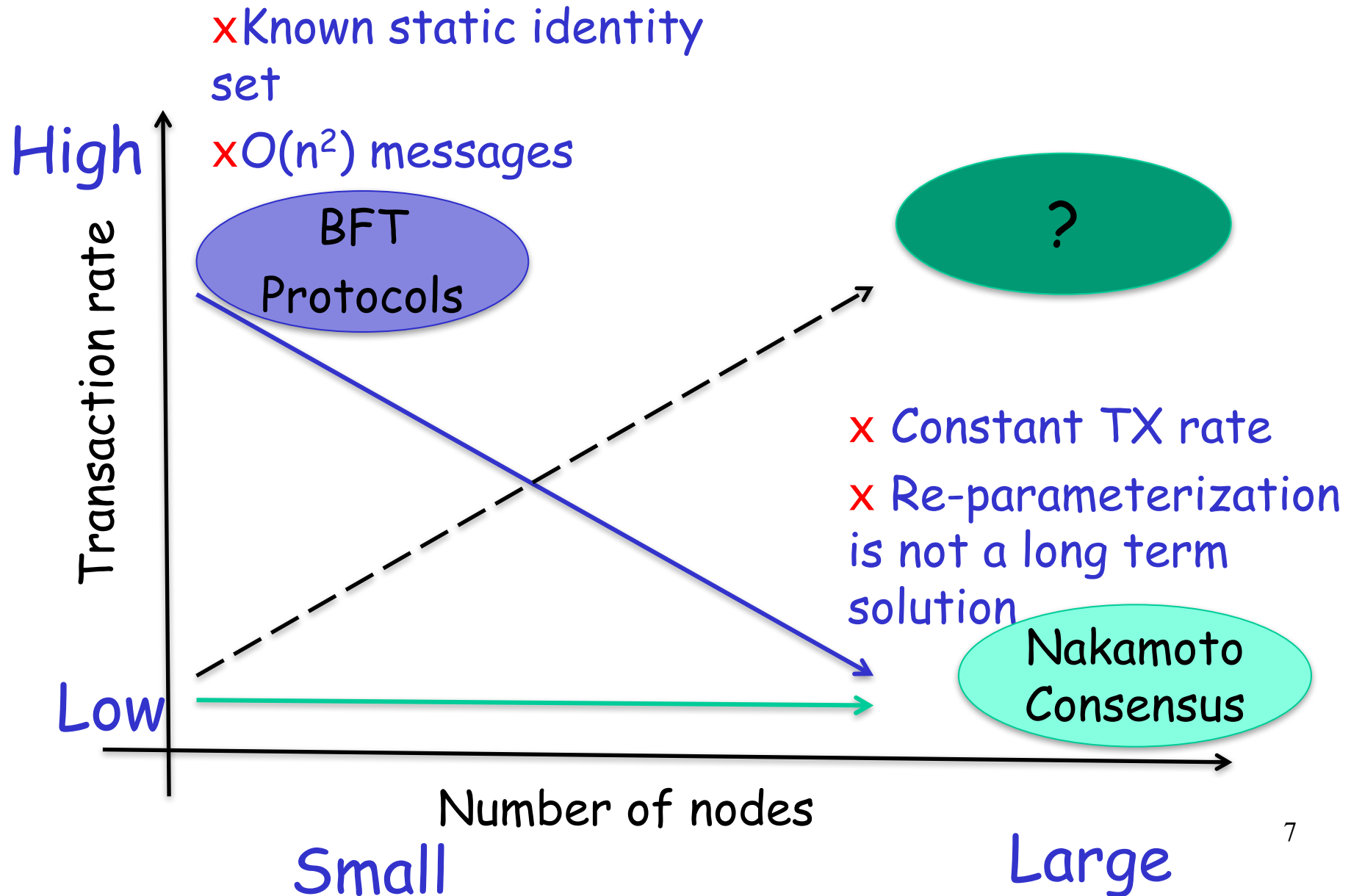
- Complete graph
- permissioned

## Agreement:

Every correct node chooses the same value

If all the correct nodes have the same input, that input must be the value chosen

# Existing protocols are not scalable



# Hybrid Consensus Idea

- Use PoW chain to elect a static committee
  - honest nodes run the blockchain for  $csize + \lambda$  blocks where  $csize = \Theta(\lambda)$  denotes the targeted committee size.
  - honest nodes would remove the trailing, unstablized  $\lambda$  blocks from its local chain, and call the miners of the first  $csize$  blocks the BFT committee.
- Run BFT to agree on a new block



# Blocks

PoW block

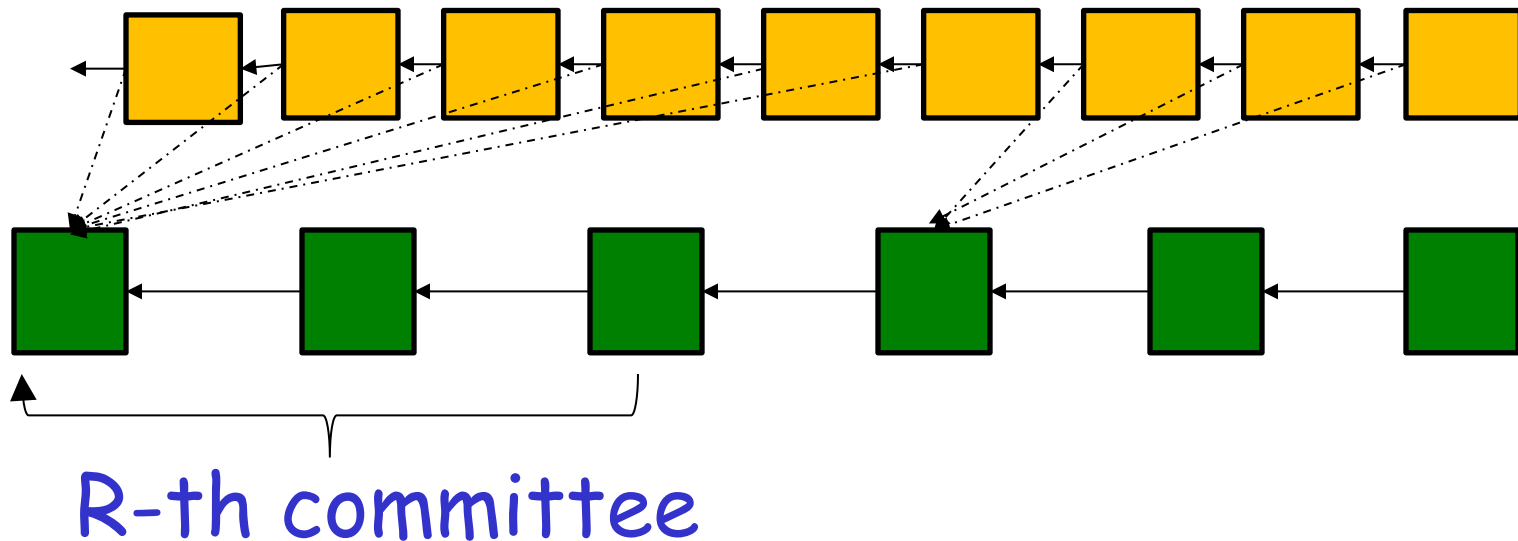
$h, \text{nonce}, \text{TXs}$

BFT block

$h', \Sigma, \text{TXs}$

$\Sigma$ : a set of signatures

$$|\Sigma| \geq \text{csize}/3$$



# Step 1: Identity establishment

## Solve PoW

$$ID = H(h, \text{Nonce}, \text{txs}) < D$$



Random seed  
for the PoW

Difficulty

The last  $csize$  confirmed blocks define  
members in a committee

## Step 2: Propose BFT blocks

Run a classical Byzantine agreement protocol

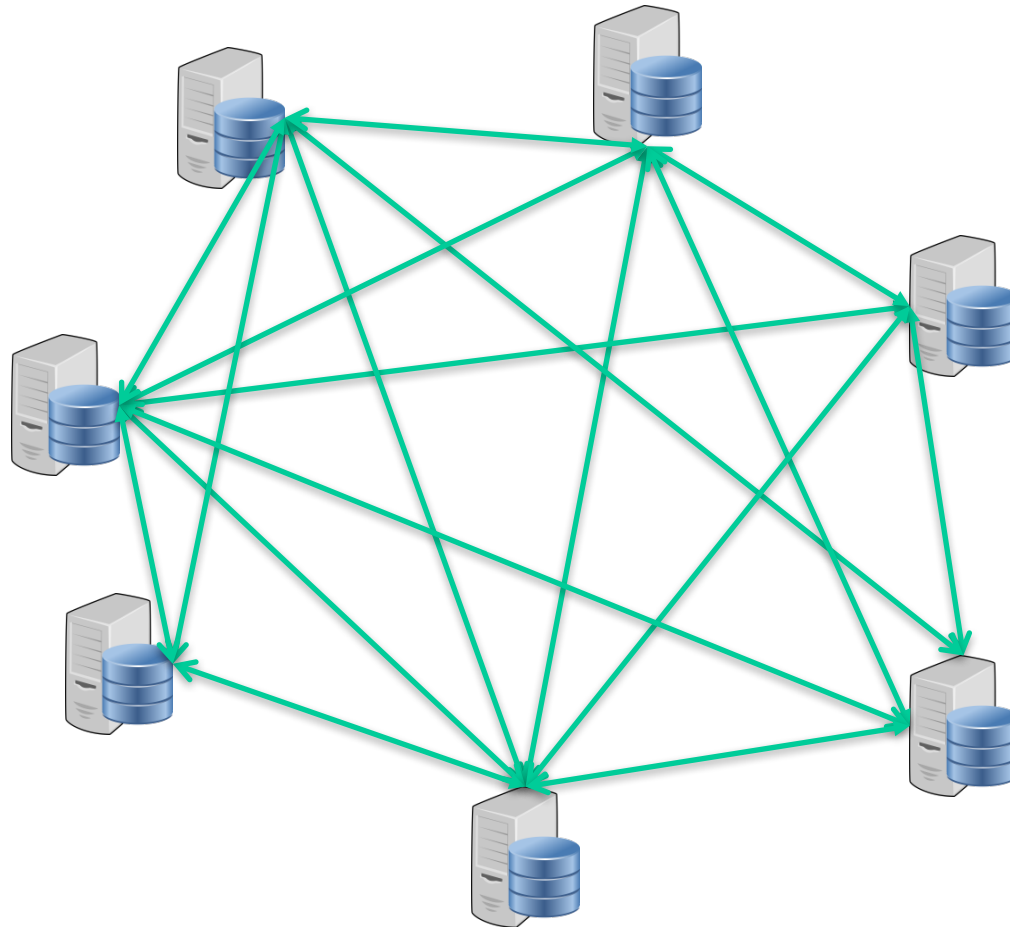
- Members agree & sign on one valid BFT block
- Each committee serves for a period of time to generate csize PoW blocks, e.g. 1 day.
- Broadcast the block to the network

When each committee has at least  $C$  members?

Each committee has  $O(\lambda)$

$O(\lambda^2)$  messages

=>scalable



# Security guarantees

Due to the consistency property of PoW chain, all honest nodes agree on the same BFT committee.

Due to the chain quality property of PoW chain, with appropriate overall parameters, we can ensure that more than  $2/3$  of the committee members are honest which is sufficient to ensure the security of the permissioned BFT protocol.

# Security guarantees

Due to the chain growth property of the pow chain, it will not take too long for the BFT committee to form

THANK YOU